

# Graph Coloring problem Using Backtracking

- Presented by:-

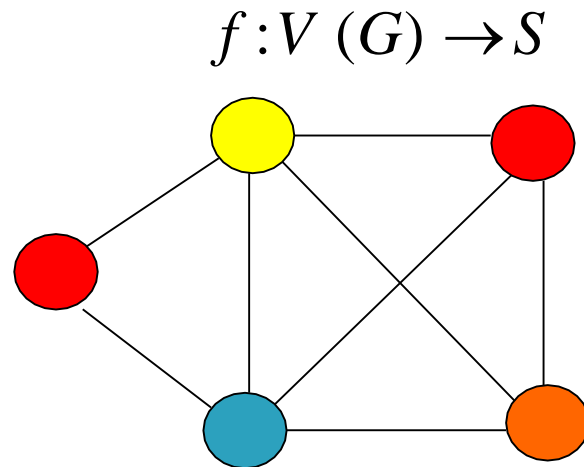
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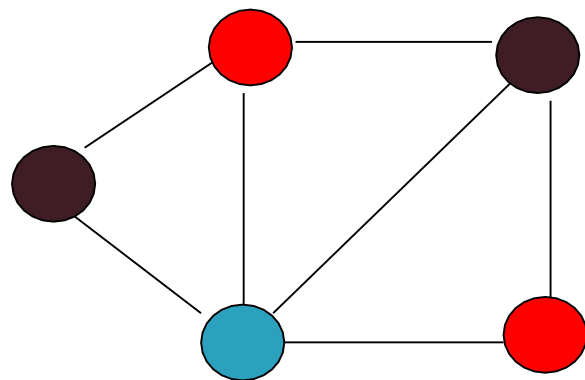
# What is Coloring?

- ▶ Graph Coloring is an assignment of colors (or any distinct marks) to the vertices of a graph. Strictly speaking, a coloring is a proper coloring if no two adjacent vertices have the same color.



# Coloring Planar graphs

- ▶ **Definition:** A graph is planar if it can be drawn in a plane without edge-crossings.



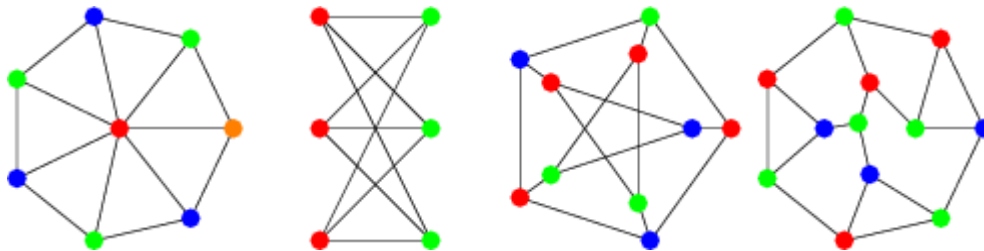
- ▶ **The four color theorem:** For every planar graph, the chromatic number is  $\leq 4$ .

# The Four-Color Theorem

- ▶ The four color theorem states that any planar map can be colored with at most four colors.
- ▶ In graph terminology, this means that using at most four colors, any planar graph can have its nodes colored such that no two adjacent nodes have the same color.
- ▶ Four-color conjecture – Francis Guthrie, 1852 (F.G.)
- ▶ Many incomplete proofs (Kempe).
- ▶ 5-color theorem proved in 1890 (Heawood)
- ▶ 4-color theorem finally proved in 1977 (Appel, Haken)
  - First major computer-based proof

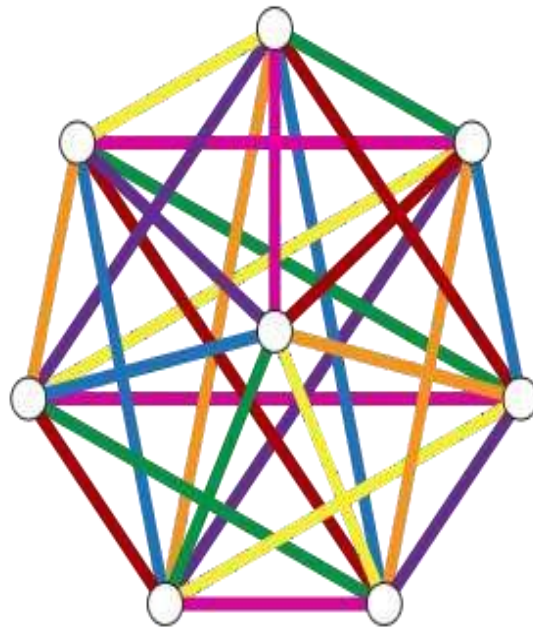
# Vertex Coloring

- ▶ A **vertex coloring** is an assignment of labels or **colors** to each **vertex** of a **graph** such that no edge connects two identically colored **vertices**

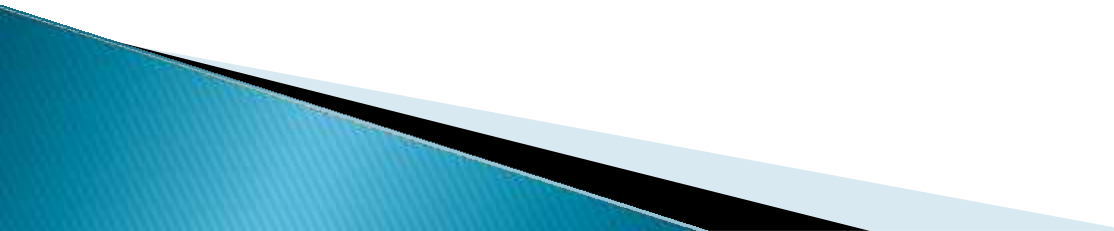


# Edge Coloring

- ▶ Similar to vertex coloring, except edges are color.
- ▶ Adjacent edges have different colors.



# Edge Coloring

- ▶ Every edge-coloring problem can be transformed into a vertex-coloring problem
  - ▶ Coloring the edges of graph  $G$  is the same as coloring the vertices in  $L(G)$
  - ▶ Not every vertex-coloring problem can be transformed to an edge-coloring problem
  - ▶ Every graph has a line graph, but not every graph is a line graph of some other graph
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# Properties

## ▶ K-Coloring

- A  $k$ -coloring of a graph  $G$  is a mapping of  $V(G)$  onto the integers  $1..k$  such that adjacent vertices map into different integers.
- A  $k$ -coloring partitions  $V(G)$  into  $k$  disjoint subsets such that vertices from different subsets have different colors.



# Terminology

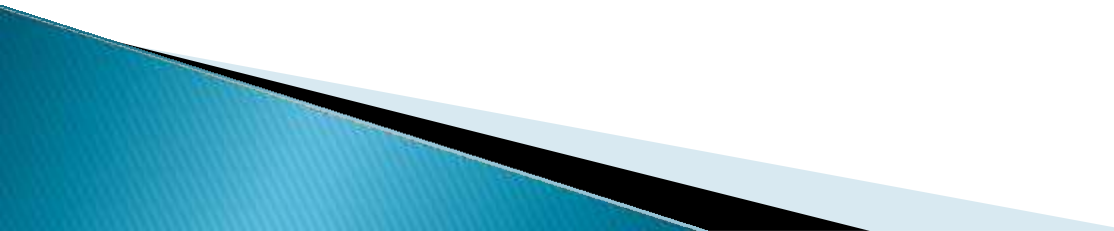
- ▶ K-colorable

- A graph  $G$  is  $k$ -colorable if it has a  $k$ -coloring.

- ▶ Chromatic Number

- The smallest integer  $k$  for which  $G$  is  $k$ -colorable is called the chromatic number of  $G$ , is denoted by the  $\chi(G)$ .

# Backtracking Algorithm

- The idea is to assign colors one by one to different vertices, starting from the vertex 0. Before assigning a color, we check for safety by considering already assigned colors to the adjacent vertices.
  - If we find a color assignment which is safe, we mark the color assignment as part of solution.
  - If we do not find a color due to clashes then we backtrack and return false.
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# NP hard problem

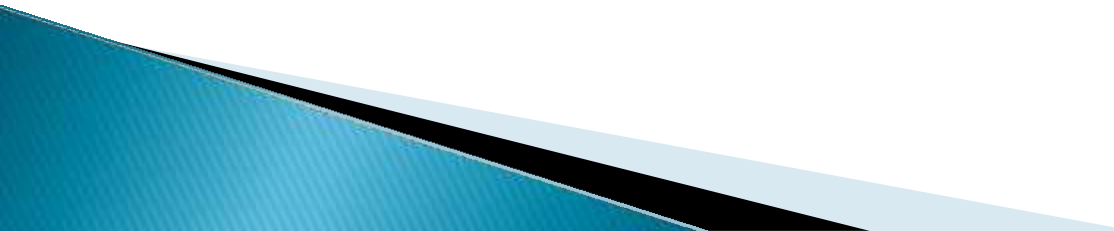
**NP**: the class of decision problems that are solvable in polynomial time on a *nondeterministic* machine (or with a nondeterministic algorithm)

- ▶ (A deterministic computer is what we know)
- ▶ A nondeterministic computer is one that can “guess” the right answer or solution think of a nondeterministic computer as a parallel machine that can freely spawn *an infinite number* of processes
- ▶ Thus *NP* can also be thought of as the class of problems whose solutions can be verified in polynomial time
- ▶ Note that *NP* stands for “Nondeterministic Polynomial-time”

# Examples Of NP hard problem

- ▶ Fractional Knapsack
- ▶ Sorting
- ▶ Others?
  - Graph Coloring
  - Satisfiability (SAT)
    - the problem of deciding whether a given Boolean formula is satisfiable.

# Applications of Graph Coloring

- ▶ Many problems can be formulated as a graph coloring problem including Time Tabling, Scheduling, Register Allocation, Channel Assignment.
  - ▶ A lot of research has been done in this area so much is already known about the problem space.
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Thank You

